

CS 477

2/6/2019

Verification using invariants in Dafny
Quantifiers

```
method m(n: nat)           := vs ==
{
    var i: int := 0;          method
    while i < n
        invariant    $\dot{i} \leq n$  assert
    {
        i := i + 1;
    }
    assert i == n;
}
```

*Initial
 $n \in \mathbb{N}$*

$$i=0 \Rightarrow i \leq n$$

method $m(n: \text{nat})$
{
var $i: \text{int} := 0;$
while $\underline{i \neq n}$
 invariant
 {
 $i := i + 1;$
 }
 assert $i == n;$
}

$i \leq n$

Initial $n \in \mathbb{N}$

$$i \leq n \wedge i \neq n$$

$$\wedge i' = i + 1$$

$$\Rightarrow i' \leq n$$

Verification Condition.

Initial $n \in \mathbb{N}$

$$i \leq n \wedge \neg(i \neq n) \Rightarrow i = n$$

```
method m(n: nat)
{
    var i: int := 2;
    while i != n
        invariant  $i \% 2 == 0$ 
    {
        i := i * 2;
    }
    assert  $i \% 2 == 0$ ;
}
```

$$R : \{i = 2^k \mid \text{for some } k\}$$



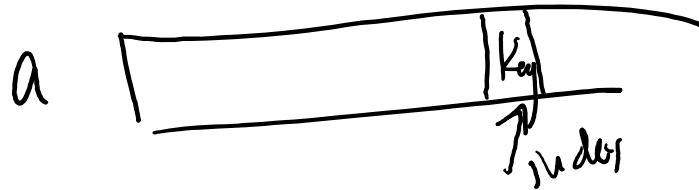
i is even

```
function fib(n: nat): nat
{
    if n == 0 then 0 else
        if n == 1 then 1 else
            fib(n - 1) + fib(n - 2)
}
method ComputeFib(n: nat) returns (b: nat)
    ensures b == fib(n)
{
    var i := 1;
    var a := 0;
    b := 1;
    while i < n
        invariant 0 < i <= n
        invariant a == fib(i - 1)
        invariant b == fib(i)
        {
            a, b := b, a + b;
            i := i + 1;
        }
}
```

```

function fib(n: nat): nat
{
    if n == 0 then 0 else
    if n == 1 then 1 else
        fib(n - 1) + fib(n - 2)
}
method ComputeFib(n: nat) returns (b: nat)
    ensures b == fib(n)
{
    if n == 0 { return 0; }
    var i := 1;
    var a := 0;
    b := 1;
    while i < n
        invariant 0 < i <= n
        invariant a == fib(i - 1)
        invariant b == fib(i)
    {
        a, b := b, a + b;
        i := i + 1;
    }
}

```



Arraya and Quantifiers

```
method Find(a: array<int>, key: int) returns (index: int)
    ensures 0 <= index ==> index < a.Length && a[index] == key
{
    index := -1
}
```

Linear Search

$$\Rightarrow \forall k \ 0 \leq k < |a| \Rightarrow a[k] \neq \text{key}$$

$$\text{index} < 0$$

```

method Find(a: array<int>, key: int) returns (index: int)
  ensures 0 <= index ==> index < a.Length && a[index] == key
  ensures index < 0 ==> forall k :: 0 <= k < a.Length ==> a[k] != key
{
  index := 0;
  while index < a.Length
    invariant forall j :: j > 0 &ampamp j < index ==> a[j] != key
    {
      if a[index] == key { return; }
      index := index + 1;
    }
  index := -1;
}

```

$$\forall j < \text{index} \quad \cancel{\forall j > 0} \quad a[j] \neq \text{key}$$

Linear Search

```
method Find(a: array<int>, key: int) returns (index: int)
    ensures 0 <= index ==> index < a.Length && a[index] == key
    ensures index < 0 ==> forall k :: 0 <= k < a.Length ==> a[k] != key
{
    index := 0;
    while index < a.Length
        invariant 0 <= index <= a.Length
        invariant forall k :: 0 <= k < index ==> a[k] != key
    {
        if a[index] == key { return; }
        index := index + 1;
    }
    index := -1;
}
```

When I get out,
index > a.length
So index = a.length.

Binary Search

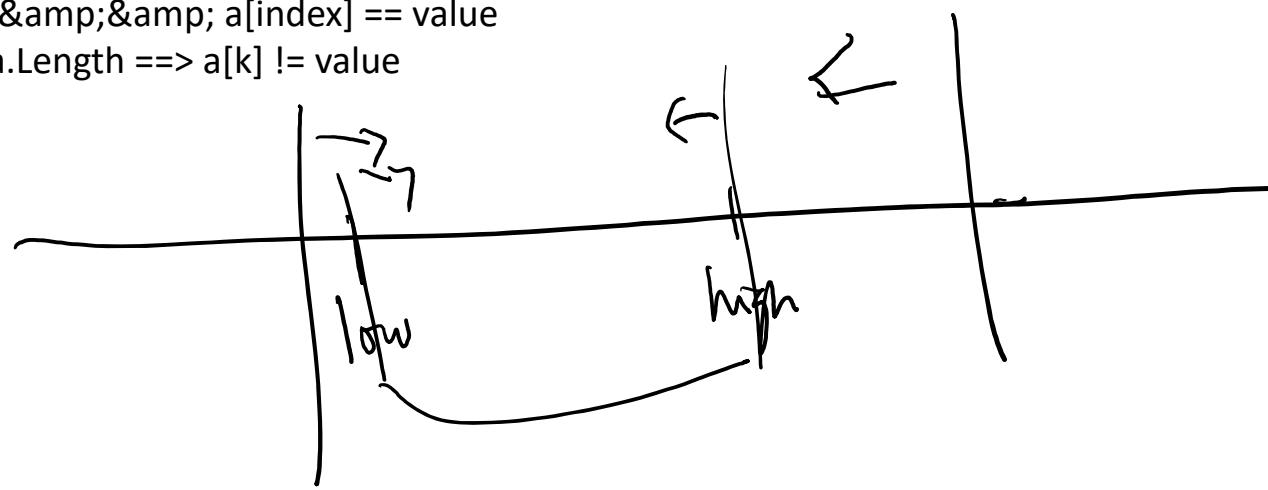
```
predicate sorted(a: array<int>)
  requires a != null
  reads a
{
  forall j, k :: 0 <= j < k < a.Length ==> a[j] <= a[k]
}
```

```
method BinarySearch(a: array<int>, value: int) returns (index: int)
  requires a != null && 0 <= a.Length && sorted(a)
  ensures 0 <= index ==> index < a.Length && a[index] == value
  ensures index < 0 ==> forall k :: 0 <= k < a.Length ==> a[k] != value
{
  // Code here
}
```

```

method BinarySearch(a: array<int>, value: int) returns (index: int)
    requires a != null && 0 <= a.Length && sorted(a)
    ensures 0 <= index ==> index < a.Length && a[index] == value
    ensures index < 0 ==> forall k :: 0 <= k < a.Length ==> a[k] != value
{
    var low, high := 0, a.Length;
    while low < high
        invariant ...
    {
        var mid := (low + high) / 2;
        if a[mid] < value
        {   low := mid + 1; }
        else if value < a[mid]
        {   high := mid; }
        else
        {   return mid; }
    }
    return -1;
}

```



44

```
method BinarySearch(a: array<int>, value: int) returns (index: int)
    requires a != null && 0 <= a.Length && sorted(a)
    ensures 0 <= index ==> index < a.Length && a[index] == value
    ensures index < 0 ==> forall k :: 0 <= k < a.Length ==> a[k] != value
{
    var low, high := 0, a.Length;
    while low < high
        invariant ...
    {
        var mid := (low + high) / 2;
        if a[mid] < value
        {
            low := mid + 1;
        }
        else if value < a[mid]
        {
            high := mid;
        }
        else
        {
            return mid;
        }
    }
    return -1;
}
```

invariant $0 \leq low \leq high \leq a.Length$
invariant $\forall i :: 0 \leq i < a.Length \wedge \neg (low \leq i < high) \Rightarrow a[i] \neq value$